

# Categorical Quandles and Knots

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# Plan

1. State Main Results
2. Motivation
3. Categories in Groups and Quandles
4. The Functors  $\text{Conj}$  and  $\text{Core}$
5. Example  $TSU(2)$
6. Non-group examples
7. local arrow systems
8. The fundamental 2-quandle

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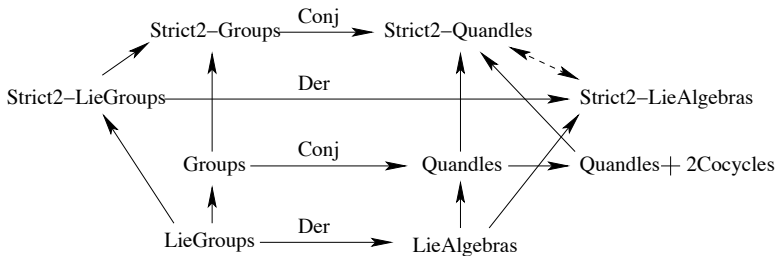
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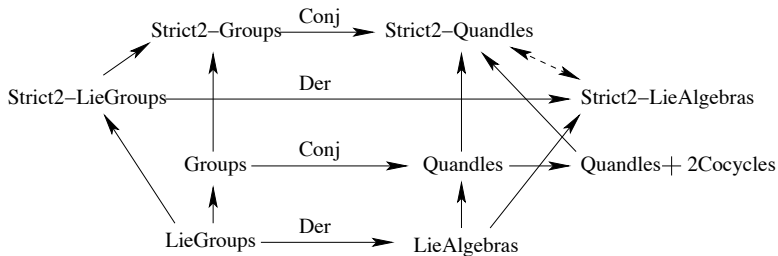
& some more stuff.

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Relations via 3-cocycles !?!?!?

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III.  $\forall x, y, z \in X \quad (x \triangleleft y) \triangleleft z = (x \triangleleft z) \triangleleft (y \triangleleft z).$

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# Category in $\mathcal{C}$

Here  $\mathcal{C}$  is either the category of groups or the category of quandles. If  $A$  is an object in  $\mathcal{C}$ , then its underlying set is denoted by  $|A|$ . A *category in  $\mathcal{C}$*  is constructed:

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 $c : (f_2, f_1) \mapsto f_2 \circ f_1$ .

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4.  $(i(x) \circ f_1) = f_1$ , and  $(f_2 \circ i(x)) = f_2$

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## Lemma

*If  $(O, M, s, t, i, \circ)_{\mathcal{C}_0}$  is a category in  $\mathcal{C}_0$ , then  $(\hat{O}, \hat{M}, \hat{s}, \hat{t}, \hat{i}, \hat{\circ})_{\mathcal{C}_1}$  is a category in  $\mathcal{C}_1$ .*

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- comp.  $((h_2, \tau(g_1)h_1) \circ (h_1, g_1) = (h_2h_1, g_1)$ .



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$$\begin{aligned}(h_1, g_1) \triangleleft (h_2, g_2) &= (\alpha(g_2^{-1}, h_2^{-1} h_1 \alpha(g_1, h_2)), g_2^{-1} g_1 g_2) \\ &= (\alpha(g_2^{-1}, h_2^{-1}) \alpha(g_2^{-1}, h_1) \alpha(g_2^{-1}, \alpha(g_1, h_2)), g_2^{-1} g_1 g_2)\end{aligned}$$

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So in the corresponding category,  $M = TSU(2)$ .

## Another gratuitous formula

$$\begin{aligned} & (a_1 \hat{i} + b_1 \hat{j} + c_1 \hat{k}, \cos(s) \hat{j} + \sin(s) \hat{k}) \\ & \triangleleft (a_2 \hat{i} + b_2 \hat{j} + c_2 \hat{k}, \cos(t) \hat{j} + \sin(t) \hat{k}) \\ & = ((2a_2 - a_1) \hat{i} \\ & \quad + [b_1 \cos(2t) \\ & + 2 \sin(s - 2t)(c_2 \cos(s) - b_2 \sin(s)) + c_1 \sin(2t)] \hat{j} \\ & \quad + [\sin(2t)(b_1 + b_2 \cos(2s) - b_2 + c_2 \sin(2s)) \\ & + \cos(2t)(-b_2 \sin(2s) - c_1 + c_2 \cos(2s) + c_2)] \hat{k}, \\ & \quad \cos(2t - s) \hat{j} + \sin(2t - s) \hat{k}) \end{aligned}$$

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- center arc next page



$$\begin{aligned}
& \left( a_2 \hat{i} + \left( \frac{-b_2 + \sqrt{3}c_2}{4} + \frac{\sqrt{3}}{2}(2c_2 - c_1) \right) \hat{j} \right. \\
& \left. - \left( \frac{1}{2}(2c_2 - c_1) + \frac{\sqrt{3}}{4}(-b_2 + \sqrt{3}c_2) \right) \hat{k}, \right. \\
& \left. \frac{-1}{2} \hat{j} - \frac{\sqrt{3}}{2} \hat{k} \right).
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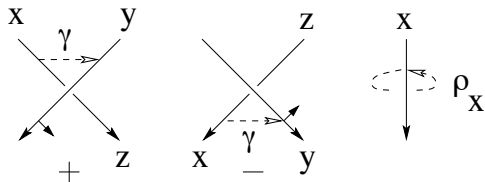
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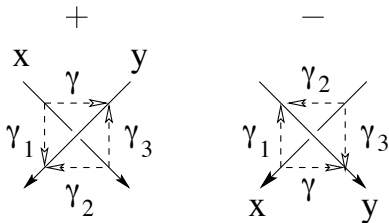
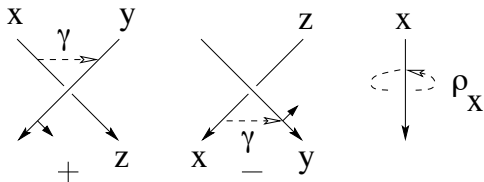
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This will be called an *Alexander 2-quandle*.

# Local Arrow Systems

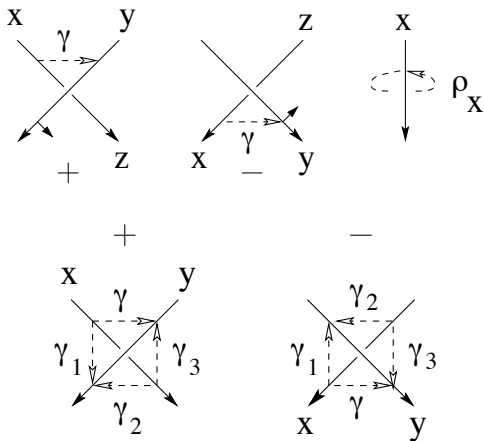


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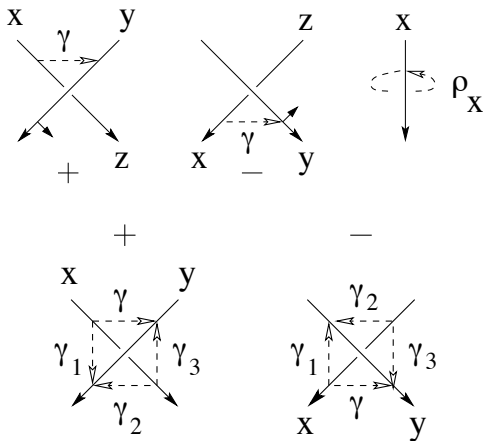


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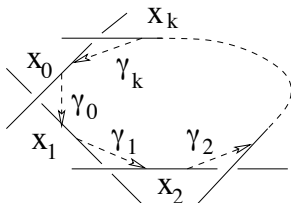


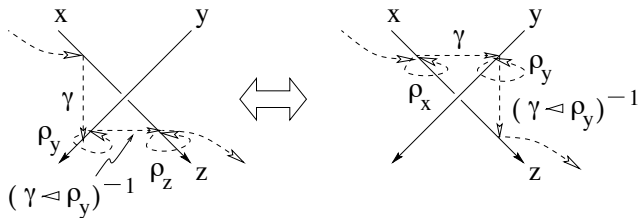
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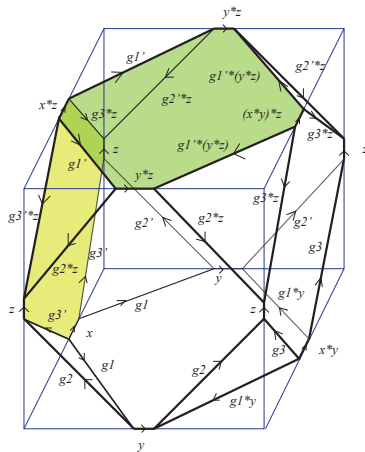




$$(\rho_z) \circ (\gamma \triangleleft \rho_y)^{-1} \circ (\rho_y) \circ (\gamma) = (\gamma \triangleleft \rho_y)^{-1} \circ (\rho_y) \circ (\gamma) \circ (\rho_x)$$

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